Exercise 2 - Cryptography

Mobile Business II (SS 2017)

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- Decrypt the following word, encrypted with the Caesar cipher:

JYFWAVNYHWOF

A	B	C	D	E	F	G	H	I	J	K	L	M
0	1	2	3	4	5	6	7	8	9	10	11	12

N	O	P	Q	R	S	T	U	V	W	X	Y	Z
13	14	15	16	17	18	19	20	21	22	23	24	25

- We assign a number for every character.
- This enables us to calculate with letters as if they were numbers.

- Encryption:
  1. Assign numbers to characters (A=0, B=1,...)
  2. Choose key k (0,..., 25)
  3. Compute $(\text{num}(\text{char}) + k) \bmod 26$, where char is the character to encrypt and $\text{num}(x)$ the number assigned to character x (e.g. $\text{num}(A) = 0$)

- How to decrypt?

- Decryption:

  1. Choose key k (0,…, 25)

  2. Assign numbers to characters (A=0, B=1,…)

  3. Compute (num(char) - k) mod 26, where char is the character to encrypt and num(x) the number assigned to character x

  4. Repeat steps for all characters

  5. Stop, if decrypted word makes sense
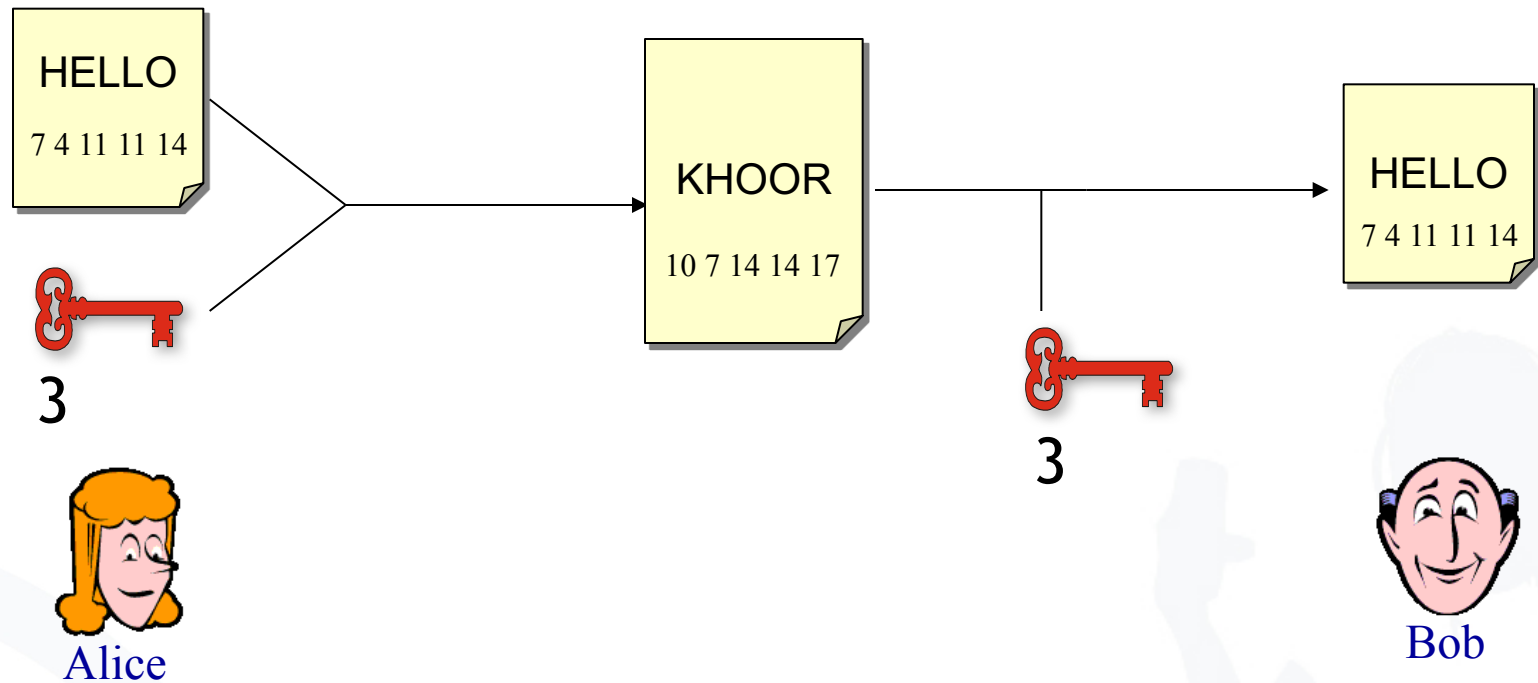
- Let's try:

Key	J	Y	F	W	A	V	N	Y	H	W	O	F
1	I	X	E	V	Z	U	M	X	G	V	N	E
2	H	W	D	U	Y	T	L	W	F	U	M	D
3	G	V	C	T	X	S	K	V	E	T	L	C
4	F	U	B	S	W	R	J	U	D	S	K	B
5	E	T	A	R	V	Q	I	T	C	R	J	A
6	D	S	Z	Q	U	P	H	S	B	Q	I	Z
7	**C**	**R**	**Y**	**P**	**T**	**O**	**G**	**R**	**A**	**P**	**H**	**Y**

Assessment of Caesar Cipher

- Very simple form of encryption.
- The encryption and decryption algorithms are very easy and fast to compute.
- It uses a very limited key space (n=26)
- Therefore, the encryption is very easy and fast to compromise.

pelcgbtencul cevbe gb gur zbqrea ntr jnf rssrpgviryl flabalzbhf jvgu rapelcgvba, gur pbairefvba bs vasbezngvba sebz n ernqnoyr fgngr gb nccnerag abafrafr. gur bevtvangbe bs na rapelcgrq zrffntr funerq gur qrpbqvat grpuavdhr arrqrq gb erpbire gur bevtvany vasbezngvba bayl jvgu vagraqrq erpvcvragf, gurerol cerpyhqvat hajnagrq crefbaf gb qb gur fnzr. fvapr jbeyq jne v naq gur nqirag bs gur pbzchgre, gur zrgubqf hfrq gb pneel bhg pelcgbybtl unir orpbzr vapernfvatyl pbzcyrk naq vgf nccyvpngvba zber jvqrfcernq.

zbqrea pelcgbtencul vf urnivyl onfrq ba zngurzngvpny gurbel naq pbzchgre fpvrapr cenpgvpr; pelcgbtencuvp nytbevguzf ner qrfvtarq nebhaq pbzchgngvbany uneqarff nffhzcgvbaf, znxvat fhpu nytbevguzf uneq gb oernx va cenpgvpr ol nal nqirefnel. vg vf gurbergvpnyyl cbffvoyr gb oernx fhpu n flfgrz ohg vg vf vasrnfvoyr gb qb fb ol nal xabja cenpgvpny zrnaf. gurfr fpurzrf ner gurersber grezrq pbzchgngvbanyyl frpher; gurbergvpny nqinaprf, r.t., vzcebirzragf va vagrtre snpgbevmngvba nytbevguzf, naq snfgre pbzchgvat grpuabybtl erdhver gurfr fbyhgvbaf gb or pbagvahnyyl nqncgrq. gurer rkvfg vasbezngvba-gurbergvpnyyl frpher fpurzrf gung cebinoyl pnaabg or oebxra rira jvgu hayvzvgrq pbzchgvat cbjre—na rknzcyr vf gur bar-gvzr cnq—ohg gurfr fpurzrf ner zber qvssvphyg gb vzcyrzrag guna gur orfg gurbergvpnyyl oernxnoyr ohg pbzchgngvbanyyl frpher zrpunavfzf.
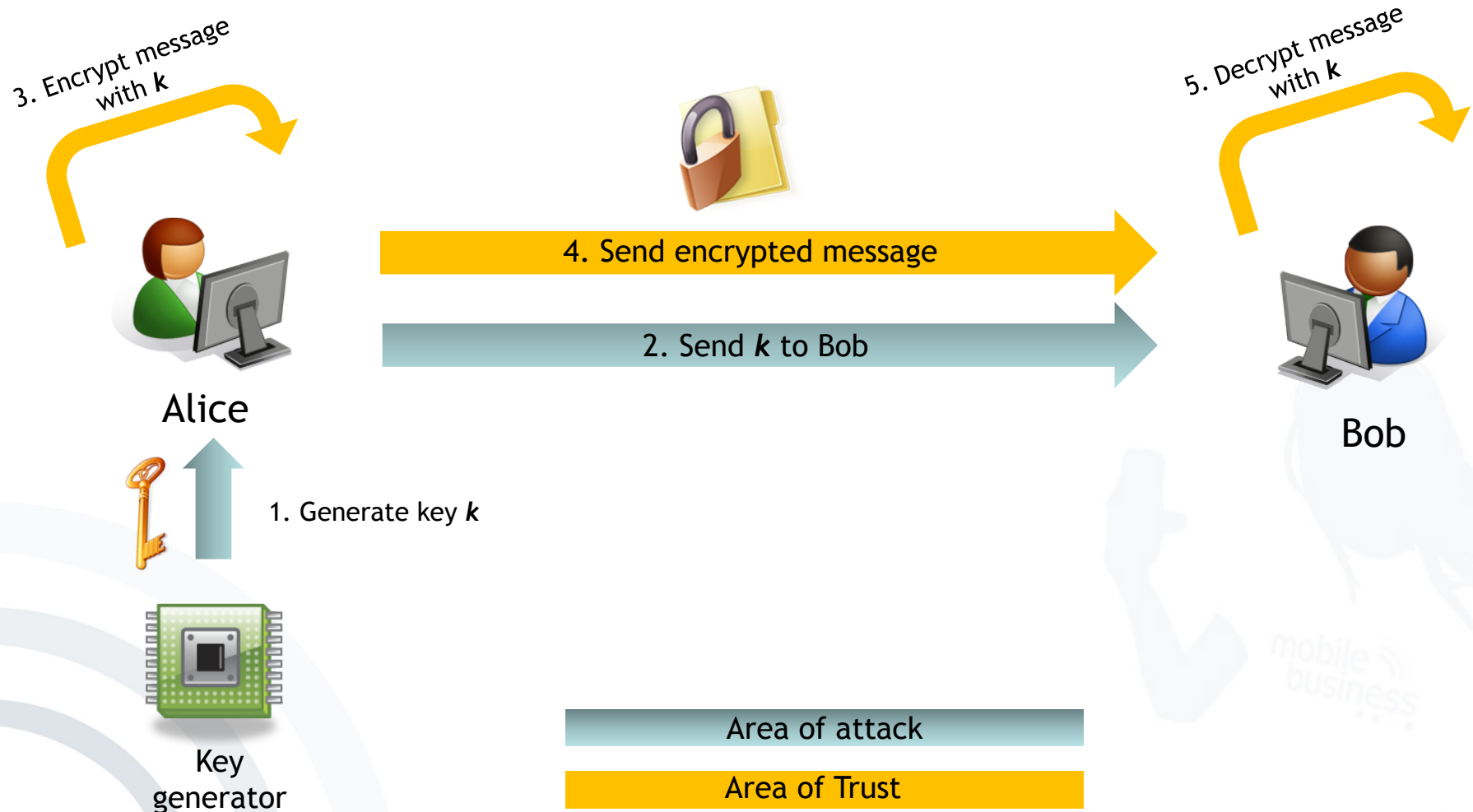
http://nayuki.eigenstate.org/page/automatic-caesar-cipher-breaker-javascript
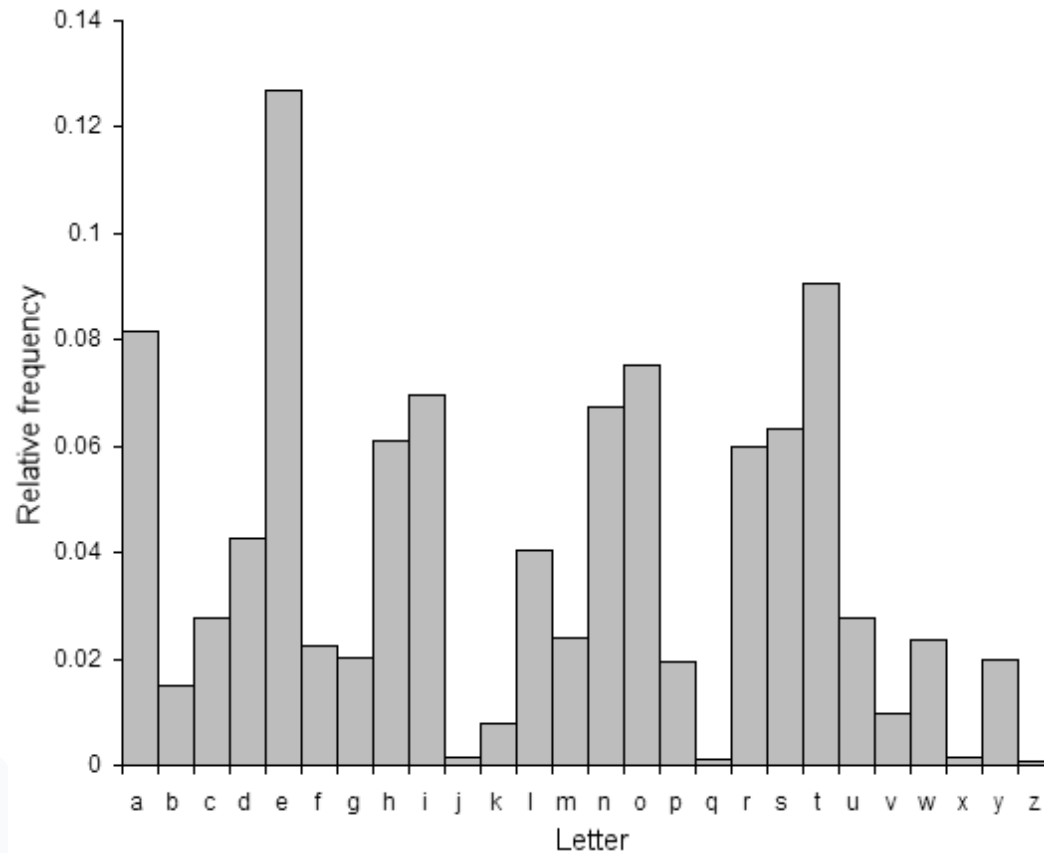
English letters frequency

1. Imagine the following situation: Alice wants to share a secret with Bob and therefore sends an encrypted message to Bob.

    1.1 Sketch the process by using symmetric encryption/decryption.
        a. Complete the illustration by highlighting each step and adding all missing elements – such as keys, involved 3rd parties,…

3. Encrypt message with *k*

5. Decrypt message with *k*

4. Send encrypted message

2. Send *k* to Bob

Alice

Bob

1. Generate key *k*

Key generator

Area of attack

Area of Trust

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b. What are pre-conditions for this approach?

b. What are pre-conditions for this approach?

→Generation of shared symmetric key

→Exchange of (secret) shared key

  →Need for secure channel

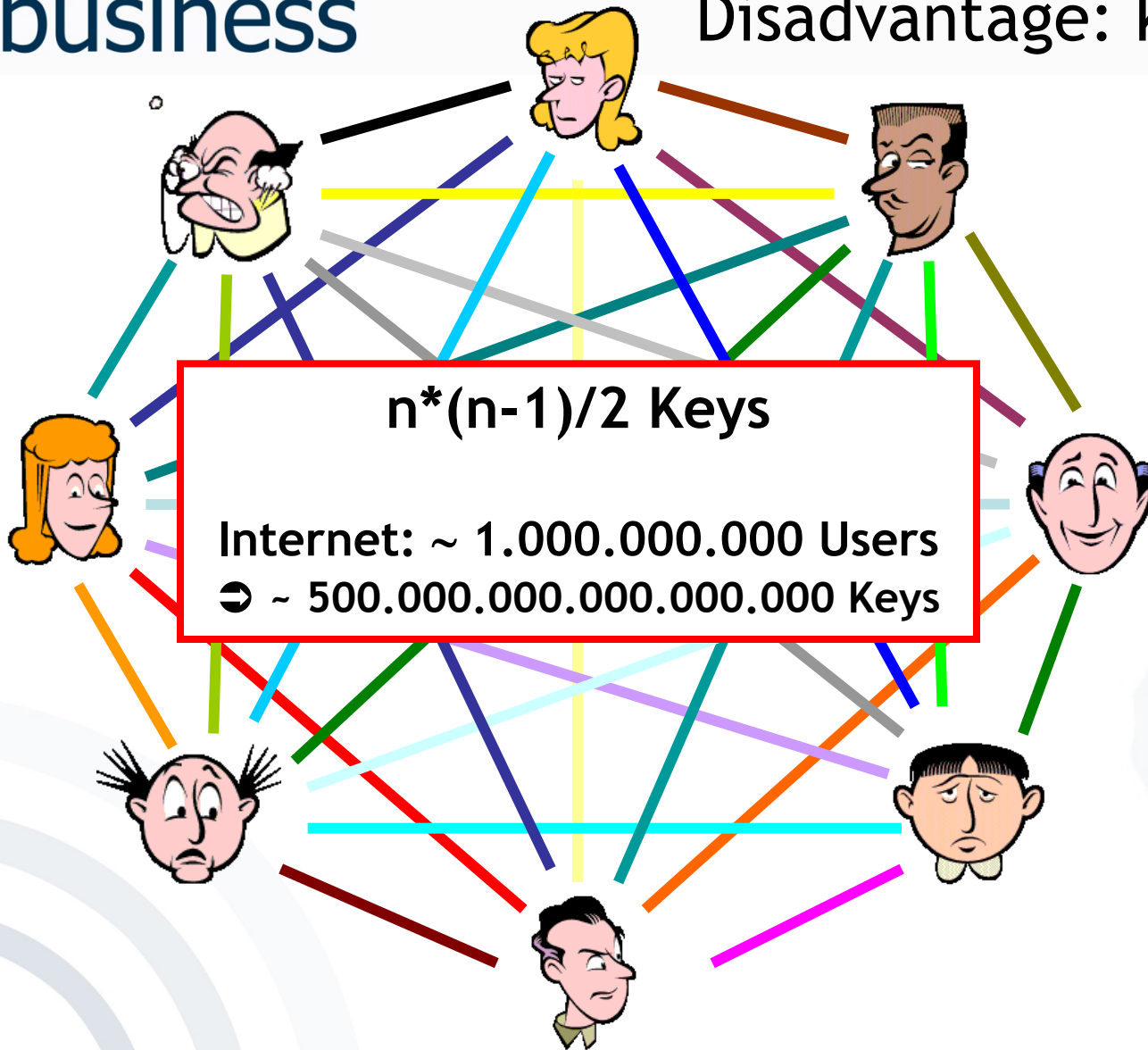c. What are advantages and disadvantages of symmetric encryption/decryption?

Advantage: Algorithms are very fast

Algorithm	Performance*
RC6	138 ms
AES	173 ms
SERPENT	200 ms
IDEA	288 ms
MARS	394 ms
TWOFISH	697 ms
DES-ede	726 ms

*) Encryption of 1 MB-blocks with an Athlon 1GHz processor

[J. Buchmann 2005: Lecture Public Key Infrastrukturen,
FG Theoretische Informatik, TU-Darmstadt]

n*(n-1)/2 Keys

Internet: ~ 1.000.000.000 Users
➲ ~ 500.000.000.000.000.000 Keys

[adopted from J. Buchmann 2005: Lecture Public Key Infrastrukturen,
FG Theoretische Informatik, TU-Darmstadt]

Key-Server

Key Server knows all secret keys!

[J. Buchmann 2005: Lecture Public Key Infrastrukturen,
FG Theoretische Informatik, TU-Darmstadt]
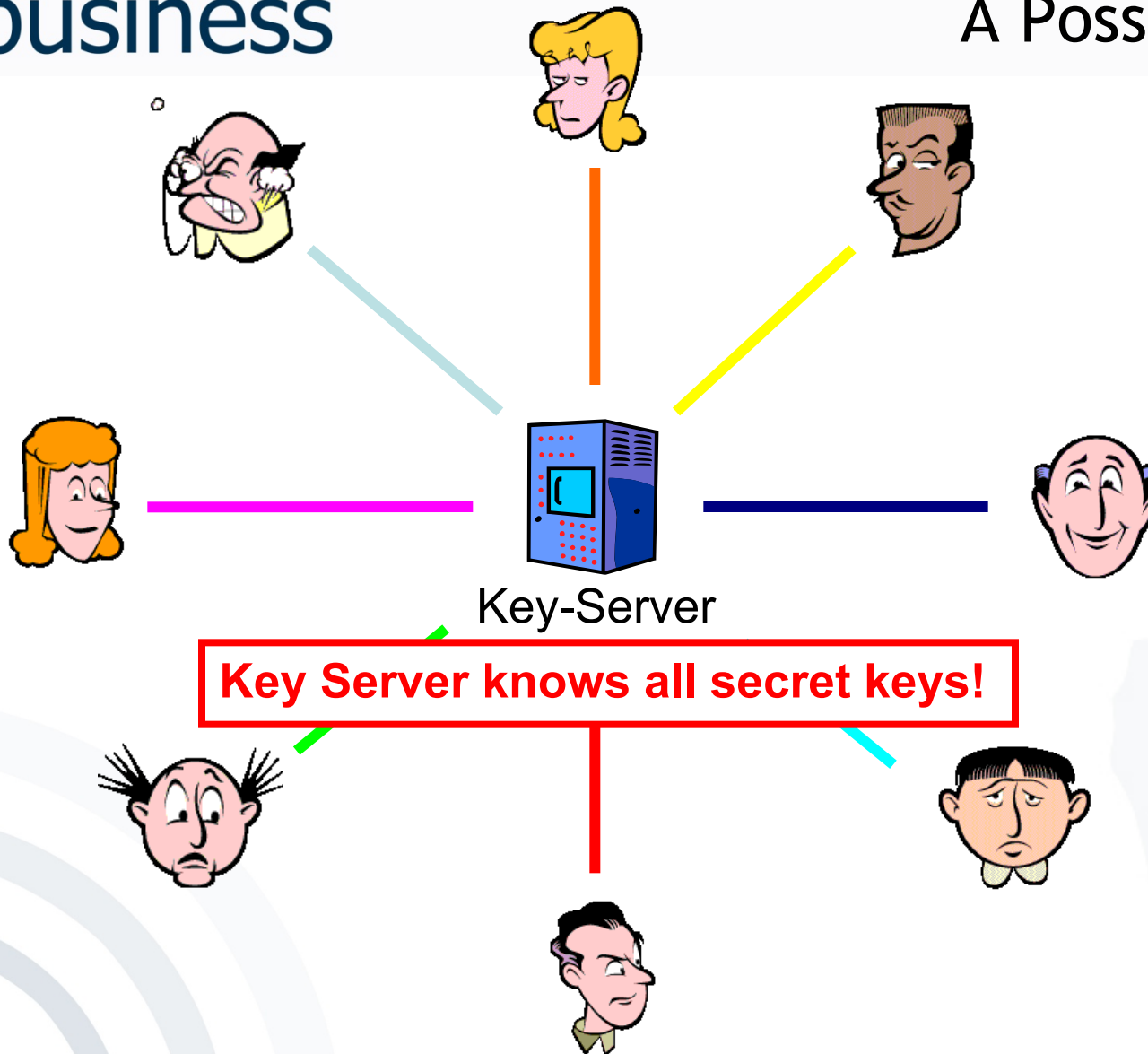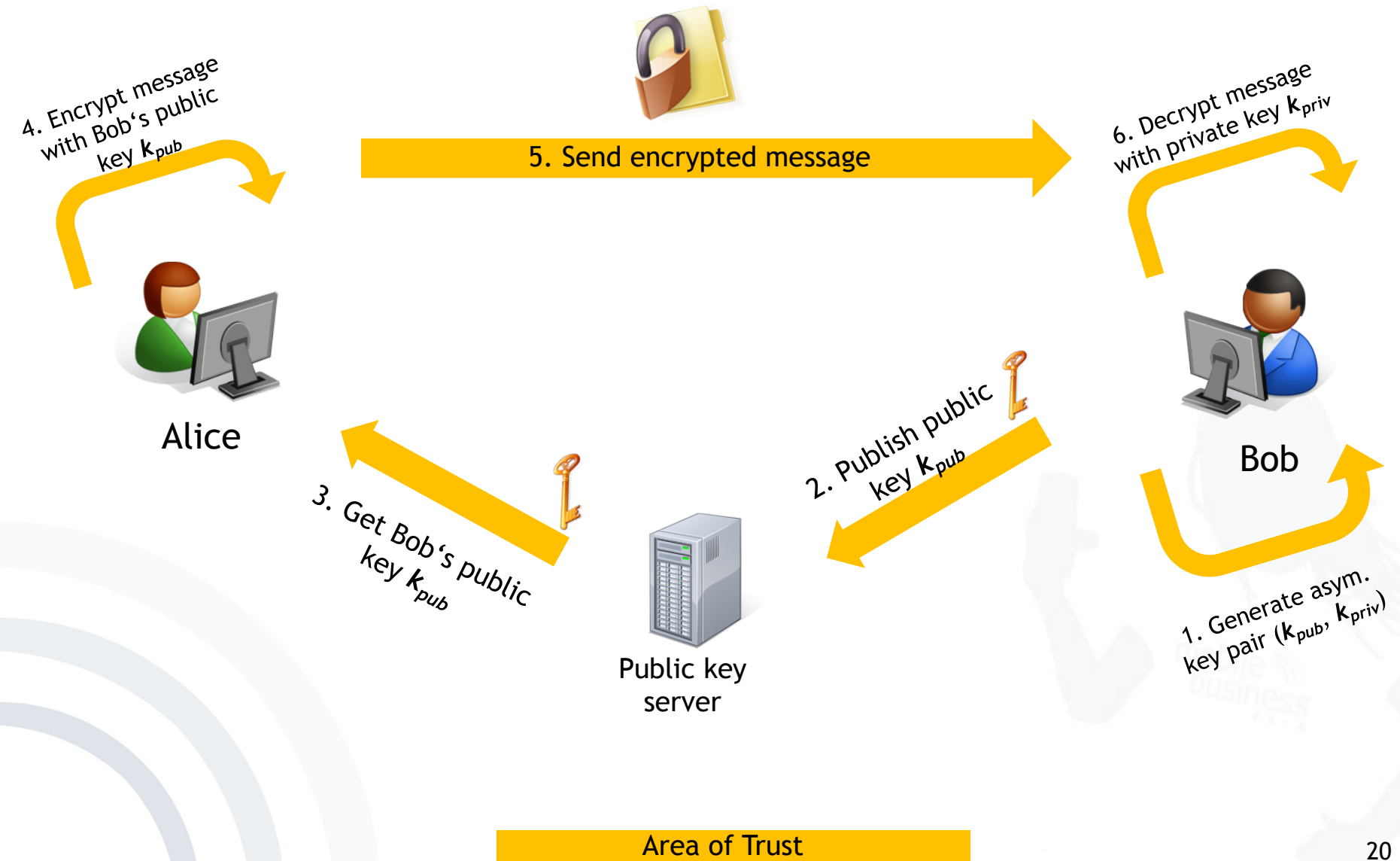
1.2 Sketch the process by using asymmetric encryption/decryption.

    a. Complete the illustration by highlighting each step and adding all missing elements – such as keys, involved 3rd parties,…

4. Encrypt message with Bob's public key k_{pub}

5. Send encrypted message

6. Decrypt message with private key k_{priv}

Alice

Bob

3. Get Bob's public key k_{pub}

2. Publish public key k_{pub}

Public key server

1. Generate asym. key pair (k_{pub}, k_{priv})

Area of Trust

b. What are pre-conditions for this approach?

b. What are pre-conditions for this approach?

→Generation of asymmetric key pairs

→Publishing public part of key

→Private key must be kept secret (!)

c. What are advantages and disadvantages of asymmetric encryption/decryption?

Algorithm	Performance*
El Gamal	1826 s
RSA	16 s

Disadvantage: Complex operations
with very big numbers

⮌ **Algorithms are very slow**

*) Encryption of 1 MB-blocks with an Athlon 1GHz processor

[J. Buchmann 2005: Lecture Public Key Infrastrukturen,
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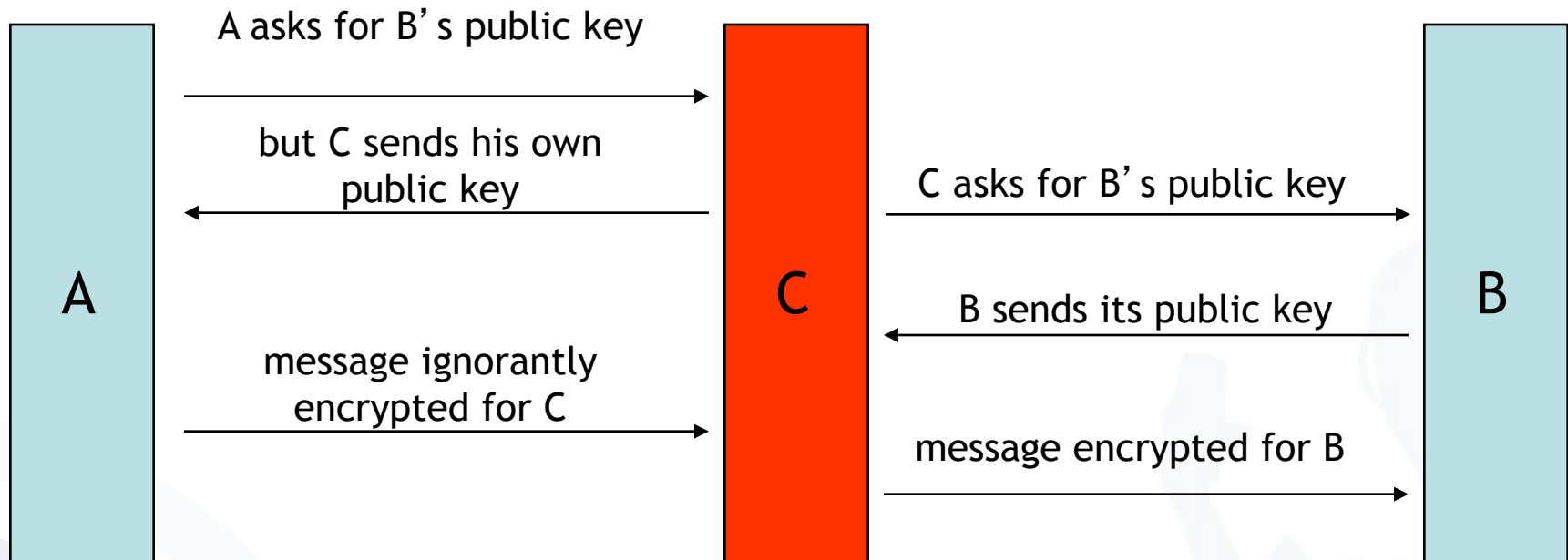c. What are advantages and disadvantages of asymmetric encryption/decryption?

Advantages:

- No secret must be shared
- Only one key per endpoint

Disadvantages:

- Algorithms are very slow
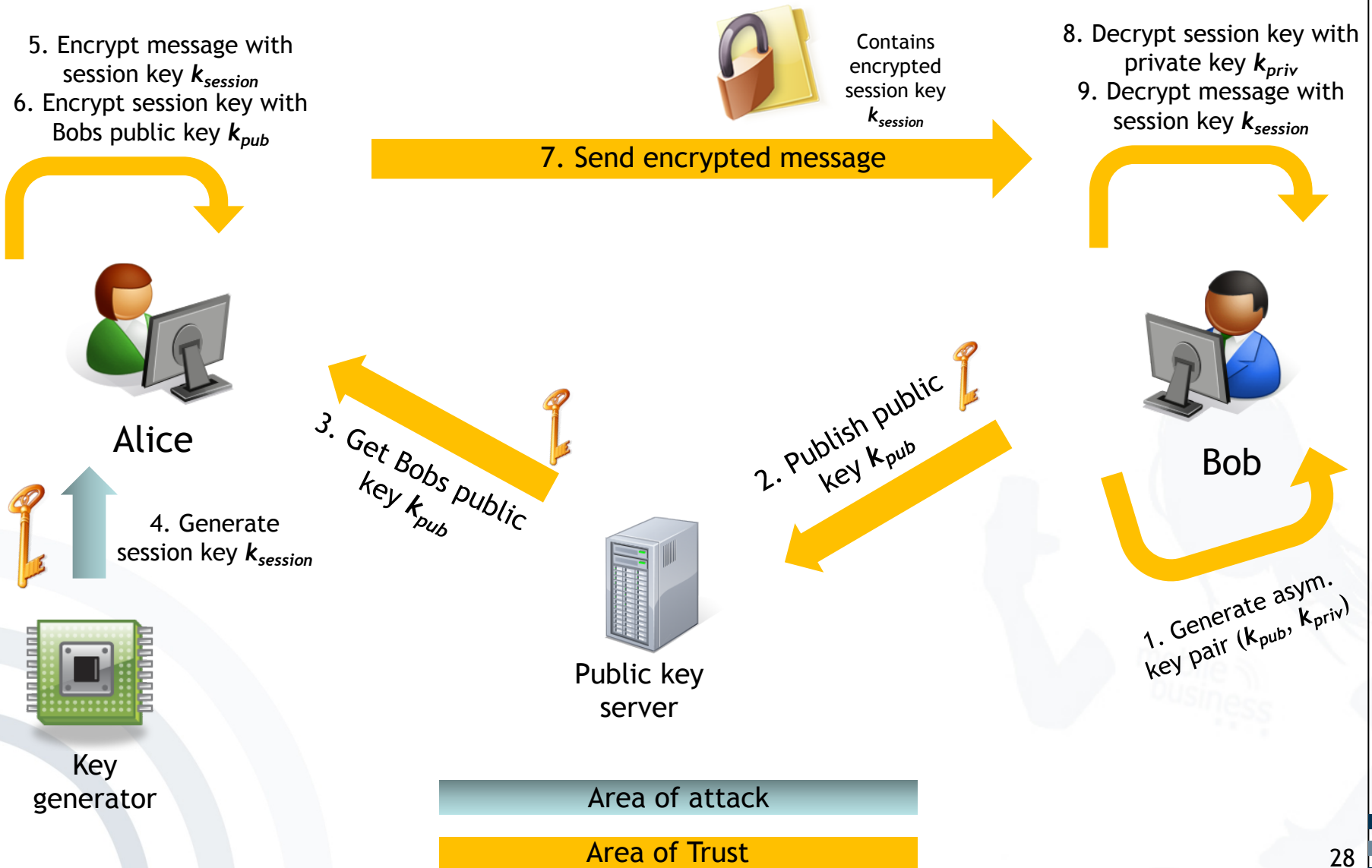- Man-in-the-middle-attack

"Man in the middle attack"

A asks for B's public key

but C sends his own public key

C asks for B's public key

A

C

B

B sends its public key

message ignorantly encrypted for C

message encrypted for B

➲ Keys are certified, that means a third person/institution confirms (with its digital signature) the affiliation of the public key to a person

1.3 Sketch the process by using PGP.
  a. Complete the illustration by highlighting each step and adding all missing elements –
     such as keys, involved 3rd parties,…

5. Encrypt message with session key $k_{session}$

6. Encrypt session key with Bobs public key k_{pub}

Contains encrypted session key $k_{session}$

8. Decrypt session key with private key k_{priv}

9. Decrypt message with session key $k_{session}$

7. Send encrypted message

Alice

3. Get Bobs public key k_{pub}

2. Publish public key k_{pub}

Bob

4. Generate session key $k_{session}$

1. Generate asym. key pair (k_{pub}, k_{priv})

Key generator

Public key server

Area of attack

Area of Trust

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b. What are pre-conditions for this approach?

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→Generation of asymmetric key pairs

→Publishing public part of key

→Private key must be kept secret (!)

→Generation of session key

c. What are advantages and disadvantages of PGP?

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→Hybrid encryption

→Advantages of both symmetric and asymmetric encryption

- **Brute-Force-Attacks on the pass phrase**
  - PGPCrack for conventionally encrypted files

- **Trojan horses, changed PGP-Code**
  - e.g. predictable random numbers, encryption with an additional key

- **Attacks on the computer of the user**
  - not physically deleted files
  - paged memory
  - keyboard monitoring

Mention possible ways for distributing keys and discuss advantages as well as disadvantages.

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➢ Manually (e.g. on flash disc)

➢ Over existing secure channel

➢ Download from (trusted) key server

➢ Stored on Smart Card

➢ Based on certificates

➢ Key exchange protocols

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